

Matching Problem

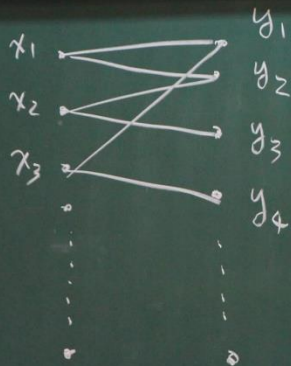
Given a bipartite graph $G = (X \cup Y, E)$,

if A is a subset of X , let

$$R(A) = \{y \in Y : \{x, y\} \in E \text{ for some } x \in A\}$$

$$R(A) = \{y \in Y : \{x, y\} \in E \text{ for some } x \in A\}$$

Example



$$A = \{x_1, x_2, x_3\}$$

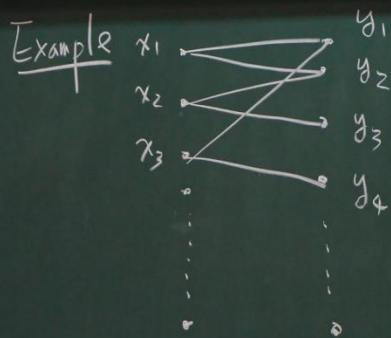
$$R(A) = \{y_1, y_2, y_3, y_4\}$$

Matching Problem

Given a bipartite graph $G = (X \cup Y, E)$,

if A is a subset of X , let

$$R(A) = \{y \in Y : \{x, y\} \in E \text{ for some } x \in A\}$$



$$A = \{x_1, x_2, x_3\}$$

$$R(A) = \{y_1, y_2, y_3, y_4\}$$

Theorem (Hall's Theorem)

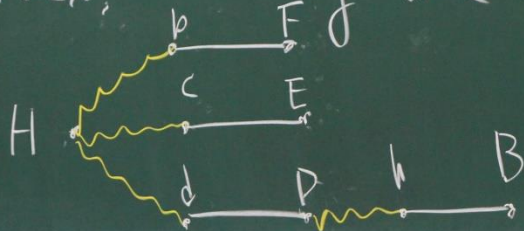
The bipartite graph $G = (X \cup Y, E)$ has a complete matching iff $|R(A)| \geq |A|$ for all $A \subseteq X$.

Proof " $\Rightarrow$ " Trivial.

" $\Leftarrow$ " If a complete matching is impossible,

the Hungarian algorithm will eventually produce a Hungarian tree, which contains one more X-vertex than Y-vertex, violating the condition.

E.g.,



Theorem (Hall's Theorem)

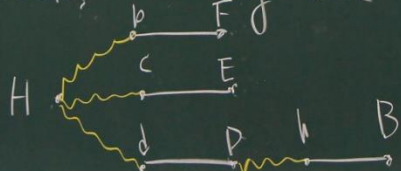
The bipartite graph $G=(X \cup Y, E)$ has a complete matching iff $|R(A)| \geq |A|$ for all $A \subseteq X$

Proof " $\Rightarrow$ " Trivial.

" $\Leftarrow$ " If a complete matching is impossible,

the Hungarian algorithm will eventually produce a Hungarian tree, which contains one more X-vertex than Y-vertex, violating the condition.

E.g.,



Def The deficiency δ of a bipartite graph $G = (X \cup Y, E)$ is defined as

$$\delta = \max_{A \subseteq X} \{|A| - |R(A)|\}.$$

Remark If $A = \phi$, then $|A| - |R(A)| = 0 - 0 = 0$.
Consequently, $\delta \geq 0$.

Def The deficiency δ of a bipartite graph $G = (X \cup Y, E)$ is defined as

$$\delta = \max_{A \subseteq X} \{|A| - |R(A)|\}.$$

Remark If $A = \phi$, then $|A| - |R(A)| = 0 - 0 = 0$.
Consequently, $\delta \geq 0$.

Property The size of a maximal matching M in a bipartite graph $G = (X \cup Y, E)$ is given by

$$|M| = |X| - \delta$$

where δ is the deficiency of G .

Property The size of a maximal matching M in a bipartite graph $G = (X \cup Y, E)$ is given by

$$|M| = |X| - \delta$$

where δ is the deficiency of G .

Def The deficiency δ of a bipartite graph $G = (X \cup Y, E)$ is defined as

$$\delta = \max_{A \subseteq X} \{|A| - |R(A)|\}$$

Remark If $A = \emptyset$, then $|A| - |R(A)| = 0 - 0 = 0$.
Consequently, $\delta \geq 0$.

Property The size of a maximal matching M in a bipartite graph $G = (X \cup Y, E)$ is given by

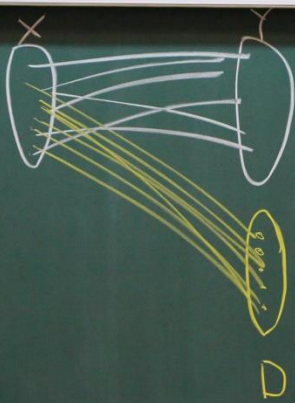
$$|M| = |X| - \delta$$

where δ is the deficiency of G .

Proof By the definition of δ , there is a subset $A_0 \subseteq X$ for which $|A_0| - |R(A_0)| = \delta$.

In any matching M , at least δ members (of A_0) remain unmatched, and $|M| \leq |X| - \delta$. We then show that there is a matching with size $|X| - \delta$.

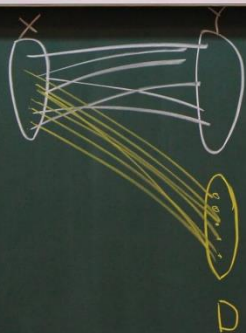
remain unmatched, and $|M| \leq |X| - \delta$. We then show that there is a matching with size $|X| - \delta$.



Let D be a new set of vertices of size δ , and let G^* be the graph $(X^* \cup Y^*, E^*)$ given by $X^* = X$, $Y^* = Y \cup D$, $E^* = E \cup K$ where K is the set of all possible edges, connecting X and D .

Proof By the definition of δ , there is a subset $A_0 \subseteq X$ for which $|A_0| - |R(A_0)| = \delta$.

In any matching M , at least δ members (of A_0) remain unmatched, and $|M| \leq |X| - \delta$. We then show that there is a matching with size $|X| - \delta$.



Let D be a new set of vertices of size δ , and let G^* be the graph $(X^* \cup Y^*, E^*)$ given by $X^* = X$, $Y^* = Y \cup D$, $E^* = E \cup K$ where K is the set of all possible edges, connecting X and D .

Hence for all $A \subseteq X (=X^*)$,

$$R^*(A) = R(A) \cup D$$

and thus

$$\begin{aligned} |R^*(A)| - |A| &= |R(A)| + |D| - |A| \\ &= \delta + \underbrace{|R(A)| - |A|}_{\geq -\delta} \geq \delta - \delta = 0 \end{aligned}$$

$$\delta = \max_{A \subseteq X} \{|A| - |R(A)|\}$$

$$\begin{aligned}
 |R^*(A)| - |A| &= |R(A)| + |D| - |A| \\
 &= \delta + \underbrace{|R(A)| - |A|}_{\geq -\delta} = \delta - \delta = 0
 \end{aligned}$$

by the definition of δ . Hence G^* satisfies Hall's condition, and it has a complete matching M^* . Removing from M^* the δ edges which have a vertex in D , we obtain the desired matching in G with size $|X| - \delta$.

Hence for all $A \subseteq X (=X^*)$,

$$R^*(A) = R(A) \cup D$$

and thus

$$\begin{aligned}
 |R^*(A)| - |A| &= |R(A)| + |D| - |A| \\
 &= \delta + \underbrace{|R(A)| - |A|}_{\geq -\delta} = \delta - \delta = 0
 \end{aligned}$$

$$\delta = \max_{A \subseteq X} \{|A| - |R(A)|\}$$

by the definition of δ . Hence G^* satisfies Hall's condition, and it has a complete matching M^* . Removing from M^* the δ edges which have a vertex in D , we obtain the desired matching in G with size $|X| - \delta$.